

COMPUTATIONAL INTELLIGENCE IN CIRCUIT SYNTHESIS THROUGH EVOLUTIONARY ALGORITHMS AND PARTICLE SWARM OPTIMIZATION

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ABSTRACT

This paper is devoted to the synthesis of combinational logic circuits through computational intelligence or, more precisely, using evolutionary computation techniques. Are studied two evolutionary algorithms, the Genetic and the Memetic Algorithm (GAs, MAs) and one swarm intelligence algorithm, the Particle Swarm Optimization (PSO). GAs are optimization and search techniques based on the principles of genetics and natural selection. MAs are evolutionary algorithms that include a stage of individual optimization as part of its search strategy, being the individual optimization in the form of a local search. The PSO is a population-based search algorithm that starts with a population of random solutions called particles. This paper presents the results for digital circuits design using the three above algorithms. The results show the statistical characteristics of this algorithms with respect to the number of generations required to achieve the solutions. The article analyzes also a new fitness function that includes an error discontinuity measure, which demonstrated to improve significantly the performance of the algorithm.

KEYWORDS: *Computational intelligence, Evolutionary algorithms, Swarm intelligence, Logic circuits design*

1. INTRODUCTION

In recent decades computational intelligence by means of evolutionary computation techniques have been applied to the design of electronic circuits and systems, leading to a novel area of research called Evolutionary Electronics (EE) or Evolvable Hardware (EH) [3]. EE considers the concept for automatic design of electronic systems. Instead of using human conceived models, abstractions and techniques, EE employs search algorithms to develop implementations not achievable with the traditional design schemes, such as the Karnaugh or the Quine-McCluskey Boolean methods [1, 4, 5].

This paper proposes three evolutionary techniques for the design of combinational logic circuits, namely a Genetic Algorithm (GA), a Memetic Algorithm (MA) and a Particle Swarm Optimization (PSO) scheme. Bearing these ideas in mind, the organization of this article is as follows. Section 2 presents the GA, the MA is described in section 3 and the PSO is detailed in section 4.

2. THE GENETIC ALGORITHM

Genetic Algorithms are adaptive heuristic search algorithms based on the evolutionary ideas of natural selection and genetic. The basic concept of GAs is designed to simulate processes in natural system necessary for evolution, specifically those that follow the principles first laid down by Charles Darwin of

survival of the fittest. As such they represent an intelligent exploitation of a random search within a defined search space to solve a problem [2]. First pioneered by John Holland in the 60s, GAs has been widely studied, experimented and applied in many fields in engineering worlds. Not only does GAs provide an alternative method to solving problem, it consistently outperforms other traditional methods in most of the problems.

GAs was introduced as a computational analogy of adaptive systems. They are modelled loosely on the principles of the evolution via natural selection, employing a population of individuals that undergo selection in the presence of operators such as mutation and recombination (crossover). A fitness function is used to evaluate individuals, and reproductive success varies with fitness.

In this section we present the adopted GA, in terms of the circuit encoding, the genetic operators and the fitness function.

2.1. Problem Definition and Circuit Encoding

A truth table specifies the circuits and the goal is to implement a functional circuit with the least possible complexity. Four sets of logic gates have been defined, as shown in **Table 1**, being Gset 2 the simplest one and Gset 6 the most complex gate set. Logic gate named WIRE means a logical no-operation.

Table 1. Gate sets.

Gate Set	Logic gates
Gset 2	{AND, XOR, WIRE}
Gset 3	{AND, OR, XOR, WIRE}
Gset 4	{AND, OR, XOR, NOT, WIRE}
Gset 6	{AND, OR, XOR, NOT, NAND, NOR, WIRE}

Section 5 exhibits the computational results. Finally, section 6 outlines the main conclusions.

In the presented scheme the circuits are encoded [6] as a rectangular matrix \mathbf{A} ($row \times column = r \times c$) of logic cells. Three genes represent each cell: $\langle input1 \rangle \langle input2 \rangle \langle gate\ type \rangle$, where $\langle input1 \rangle$ and $\langle input2 \rangle$ are one of the circuit inputs, if they are in the first column, or one of the previous outputs, if they are in other columns. The gate type is one of the elements adopted in the gate set. The chromosome is formed with as many triplets as the matrix size demands (e.g. triplets = $3 \times r \times c$).

2.2. The Genetic Operators

The initial population of circuits (strings) is generated at random. The search is then carried out among this population. The three different operators used are reproduction, crossover and mutation, as described in the sequel. In what concern the reproduction operator, the successive generations of new strings are reproduced on the basis of their fitness function. In this case, it is used a tournament selection to select the strings from the old population, up to the new population. For the crossover operator, the strings in the new population are grouped together into pairs at random. Single point crossover is then performed among pairs. The crossover point is only allowed between cells to maintain the chromosome integrity.

The mutation operator changes the characteristics of a given cell in the matrix. Therefore, it modifies the gate type and the two inputs, meaning that a completely new cell can appear in the chromosome. Moreover, it is applied an elitist algorithm and, consequently, the best solutions are always kept for the next generation. To run the GA we have to define the number of individuals to create the initial population P. This population is always the same size across the generations, until the solution is reached. The crossover rate CR represents the percentage of the population P that reproduces in each generation. Like-wise, the mutation rate MR is the percentage of the population P that can mutate in each generation.

2.3. THE FITNESS FUNCTION

The calculation of the fitness function F in Eq. (1) has two parts, f_1 and f_2 , where f_1 measures the functionality and f_2 measures the simplicity. In a first phase, we compare the output \mathbf{Y} produced by the GA-generated circuit with the required values \mathbf{Y}_R , according with the truth table, on a bit-per-bit basis. By other words, f_1 is incremented by one for each correct bit of the output until f_1 reaches the maximum value f_{10} that occurs when we have a functional circuit. Once the circuit is functional, in a Second phase, the algorithm tries to generate circuits with the least number of gates. This means that the resulting circuit must have as much genes $gate\ type \equiv wire$ as possible. Therefore, the index f_2 , that measures the simplicity (the number of null operations), is increased by *one (zero)* for each *wire (gate)* of the generated circuit, yielding:

• **First phase, circuit functionality:**

$$f_{10} = 2^{ni} \times no \dots \dots \dots (1a)$$

$$f_{11} = f_{11} + 1,$$

$$if \{bit\ i\ of\ \mathbf{Y}\} = \{bit\ i\ of\ \mathbf{Y}_R\}$$

$$i = 1, \dots, f_{10} \dots \dots \dots (1b)$$

$$f_1 = f_{11} - \delta,$$

$$if\ error_i \neq error_{i-1},\ i = 1, \dots, f_{10} \dots \dots \dots (1c)$$

(When measuring discontinuity)

• **Second phase, circuit simplicity:**

$$f_2 = f_2 + 1, \text{ if } gate\ type = wire \dots \dots \dots (1d)$$

$$F = \begin{cases} f_1, & F < f_{10} \\ f_1 + f_2, & F \geq f_{10} \end{cases} \dots \dots \dots (1e)$$

where $i = 1, \dots, f_{10}$, ni and no represent the number of inputs and outputs of the circuit.

3. THE MEMETIC ALGORITHM

This section describes the MA. MAs are inspired by models of adaptation in natural systems that combine evolutionary adaptation of populations of individuals with individual learning within a lifetime [10]. As it is known, MAs are Meta heuristics that take advantage of the evolutionary operators in determining interesting regions of the search space. Moreover, MAs adopt a local search that rapidly finds good solutions in a small region of the search space [11]. Additionally, MAs are inspired by Richard Dawkins' concept of a meme, which represents a unit of cultural evolution that can exhibit local refinement [12]. The proposed MA includes a GA and a local search algorithm, where the GA corresponds to the algorithm implemented in first stage of development. Over the last decade, MAs have relied on the use of a variety of different methods as the local improvement procedure. Some recent studies on the choice of local search method employed have shown that this choice significantly affects the efficiency of problem searches [13].

The local search method investigates a small area around a solution and adopts the best-found solution. By other words, the procedure tries to find a fitter solution in the neighborhood of the current solution. If the algorithm finds a better solution, then the new solution replaces the current solution, and the neighborhood restarts. Local search methods are iterative algorithms that seek to enhance the solution by stepwise improvements. The simplest form of local search attempts to swap elements in combinatorial optimization problems. In our case, it is implemented a gate type local search (GTLS) algorithm as shown in Fig. 1.

4. THE PARTICLE SWARM OPTIMIZATION

Particle Swarm Optimization (PSO) is a recently proposed algorithm by James Kennedy and R. C. Eberhart. in 1995, motivated by social behavior of organisms such as bird flocking and fish schooling.

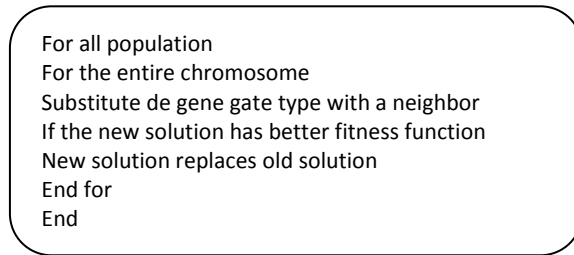


Fig. 1. The local search algorithm.

PSO algorithm is not only a tool for optimization, but also a tool for representing sociocognition of human and artificial agents, based on principles of social psychology [14]. Some scientists suggest that knowledge is optimized by social interaction and thinking is not only private but also interpersonal. PSO as an optimization tool, provides a population-based search procedure in which individuals called particles change their position (state) with time. In a PSO system, particles fly around in a multidimensional search space. During flight, each particle adjusts its position according to its own experience, and according to the experience of a neighboring particle, making use of the best position encountered by itself and its neighbor. Thus, as in modern GAs and MAs, a PSO system combines local search methods with global search methods, attempting to balance exploration and exploitation [15].

4.1. PSO Parameters

In the PSO, instead of using genetic operators, as in the case of GAs, each particle (individual) adjusts its flying according with its own and its companion’s experiences. Each particle is treated as a point in a D-dimensional space and is manipulated as described below in the original PSO algorithm:

$$v_{id} = v_{id} + c_1 \text{rand}() (p_{id} - x_{id}) + c_2 \text{rand}() (p_{gd} - x_{id}) \dots \dots \dots (2)$$

$$x_{id} = x_{id} + v_{id} \dots \dots \dots (3)$$

Where c_1 and c_2 are positive constants and $\text{rand}()$ is a random function in the range $[0,1]$, $X_i = (x_{i1}, x_{i2}, \dots, x_{iD})$ represents the i th particle, $P_i = (p_{i1}, p_{i2}, \dots, p_{iD})$ is the best previous position (the position giving the best fitness value) of the particle, the symbol g represents the index of the best particle among all particles in the population, and $V_i = (v_{i1}, v_{i2}, \dots, v_{iD})$ is the rate of the position change (velocity) for particle i . Expression (2) represents the flying trajectory of a population of particles. Eq. (2) describes how the velocity is dynamically updated and Eq. (3) the position update of the “flying” particles. Eq. (2) is divided in three parts, namely the momentum, the cognitive and the social parts. In the first part the velocity cannot be changed abruptly: it is adjusted based on the current velocity. The second part represents the learning from its own flying experience. The third part consists on the learning group flying experience [15, 16].

The initial velocity of each particle is initialized with zero. The velocities of the following generations are calculated applying Eq. (2) and the new positions result from using Eq. (3). In this way, each potential solution, called particle, flies through the problem space. For each gene is calculated the corresponding velocity. Therefore, the new positions are as many as the number of genes in the chromosome. If the new values of the input genes result out of range, then a re-insertion function is used. If the calculated gate gene is not allowed a new valid one is generated at random. These particles then have memory and each one keeps information of its previous best position ($pbest$) and its corresponding fitness. The swarm has the $gbest$ of all the particles and the particle with the greatest fitness is called the global best ($gbest$). The basic concept of the PSO technique lies in accelerating each particle towards its $pbest$ and $gbest$ locations with a random weighted acceleration. However, in our case we also use a kind of mutation operator that

introduces a new cell in 10% of the population. This mutation operator changes the characteristics of a given cell in the matrix. Therefore, the mutation modifies the gate type and the two inputs, meaning that a completely new cell can appear in the chromosome. To run the PSO we have also to define the number P of individuals to create the initial population of particles. This population is always the same size across the generations, until reaching the solution.

5. COMPUTATIONAL RESULTS

This section presents the computational results when applying, firstly the evaluation of the fitness function without the error discontinuity measure and secondly using the fitness function with the error discontinuity measure. Due to the stochastic nature of the EAs, in order to evaluate its performance, for each gate set we perform 20 simulations. The best gate set is the one that requires the smaller average number of generations $Av(N)$ and the smaller standard deviation $S(N)$ to reach the solution.

5.1. Using the Fitness without Discontinuity Measure

This subsection shows the implementation of four different combinational logic circuits, namely, a 2-to-1 multiplexer (M2-1), a one-bit full adder (FA1),

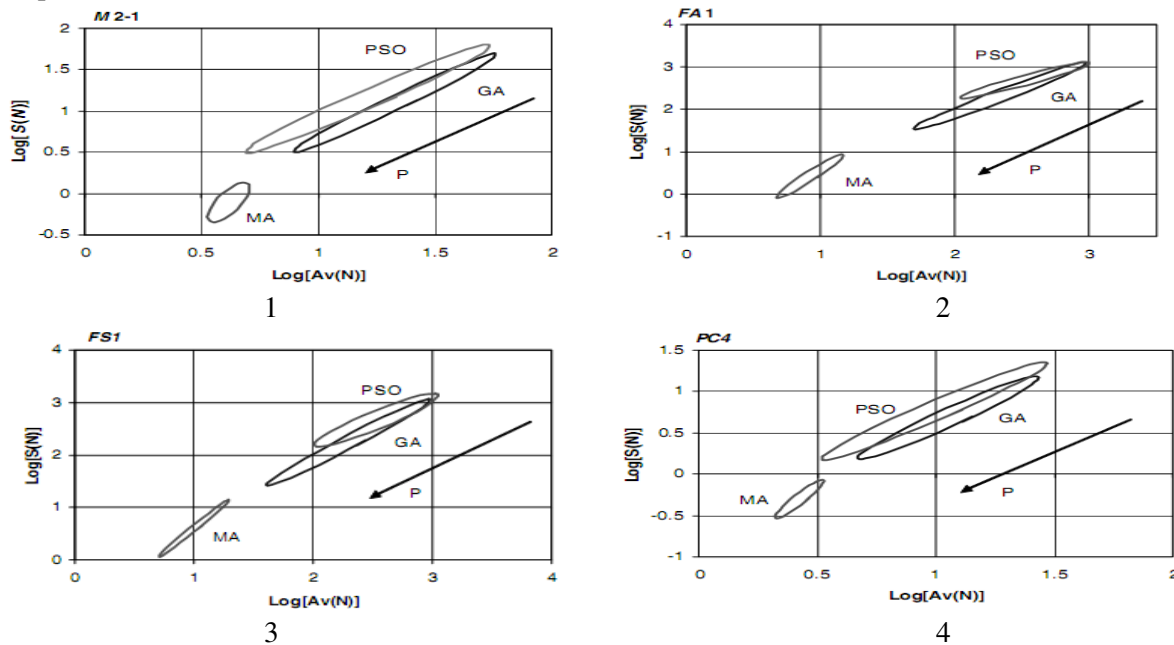


Fig. 2. $\text{Log}[S(N)]$ versus $\text{Log}[Av(N)]$ for the M2-1, the FA1, the FS1 and the PC4 circuits with $P=\{100, 500, 1000, 3000\}$ and $\delta=0$.

a one-bit full subtractor (FS1) and a four-bit parity checker (PC4), using the GA, the MA and the PSO algorithms, using the fitness function described above, without the discontinuity error measure, that is with $\delta=0$.

The first case study is the M2-1 circuit, with a truth table with three inputs $\{S_0, I_1, I_0\}$ and one output $\{O\}$. The matrix has a size of $r \times c = 3 \times 3$ and the length of each string representing a circuit (i.e., the chromosome length) is $CL=27$. Since the 2-to-1 multiplexer has $n_i=3$ and $n_o=1$, it results $f_{10}=8$ and $F \geq 12$.

The second case study is the FA1 circuit, with a truth table with three inputs $\{A, B, C_{in}\}$ and two outputs $\{S, C_{out}\}$. In this case, the matrix has a size of $r \times c = 3 \times 3$, and the chromosome length is $CL=27$. Since the one-bit full adder has $n_i=3$ and $n_o=2$, it results $f_{10}=16$ and $F \geq 20$.

The third case study is a FS1 circuit, with a truth table with three inputs $\{A, B, B_{in}\}$ and two outputs $\{S, B_{out}\}$. In this case, the matrix has a size of $r \times c = 3 \times 3$, and the chromosome length is $CL=27$. Since the one-bit full subtractor has $n_i=3$ and $n_o=2$, it results $f_{10}=16$ and $F \geq 20$.

The fourth case study consists on the PC4 circuit, which has four inputs $\{A_3, A_2, A_1, A_0\}$ and one output $\{O\}$. The size of the matrix is $r \times c = 4 \times 4$ and the chromosome length is $CL=48$. In this case $n_i=4$ and $n_o=1$, resulting $f_{10} = 16$ and $F \geq 24$.

Figure 2 presents the results obtained in terms of $\text{Log}[S(N)]$ versus $\text{Log}[Av(N)]$ for the M2-1, the FA1, the FS1 and the PC4 circuits and $P = \{100, 500, 1000, 3000\}$ with $\delta = 0$.

The points in the space $\{\text{Log}[Av(N)], \text{Log}[S(N)]\}$ are Approximated by a bi-dimensional Gaussian probability distribution. The ellipses depicted in the charts represent the corresponding contour plots.

It is obvious that the MA algorithm reveals a better performance for all the combinational circuits and that both $Av(N)$ and $S(N)$ vary inversely with P . The GA and the PSO algorithms present similar results in particular for the M2-1 and the PC4 circuits. For the FA1 and the FS1 the PSO is less sensitive to P than the GA.

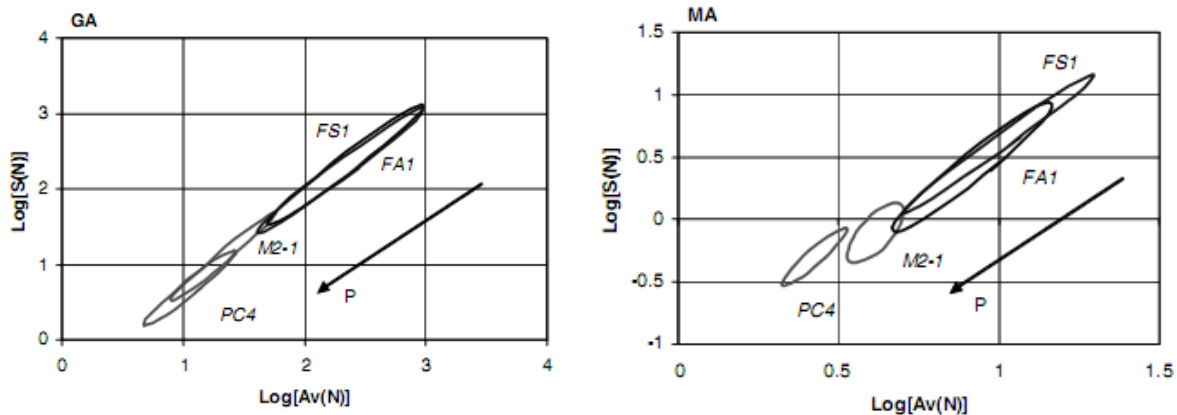
Figure 3 shows $\text{Log}[S(N)]$ versus $\text{Log}[Av(N)]$ with $P = \{100, 500, 1000, 3000\}$ for the GA, the MA and the PSO algorithms, with $\delta = 0$.

Analyzing the charts is possible to classify the complexity of the combinational logic circuits in the perspective of each evolutionary algorithm. For the three algorithms, the sequence of increasing circuit complexity becomes $\{PC_4, M_2-1, FA_1, FS_1\}$. In the PSO algorithm, the circuit complexity is clearly divided in two zones, namely the $\{FS_1, FA_1\}$ and the $\{M_2-1, PC_4\}$ zones.

5.2. Using the Fitness with Discontinuity Measure

The experiments of this subsection consist on running the PSO algorithm to generate the one-bit full adder (FA1) using the fitness scheme described in Eq. (1). The circuit is generated with the gate sets presented in Table 1 and $P = 3000$, $w = 0.5$, $c_1 = 1.5$ and $c_2 = 2$.

Figure 4 depict the average number of generations $Av(N)$ and the standard deviation of the number of generations $S(N)$ to achieve the solution versus δ for the PSO algorithm, the circuit $\{FA1\}$ and the gate sets $\{2, 3, 4, 6\}$. The results reveal that *Gset 3* presents a superior performance to the other *Gsets*, for all values of δ . Moreover, analyzing the influence of δ we conclude that the PSO response is better mostly in the region $\delta \approx 0.5$ for the arithmetic circuit and for all gate sets.



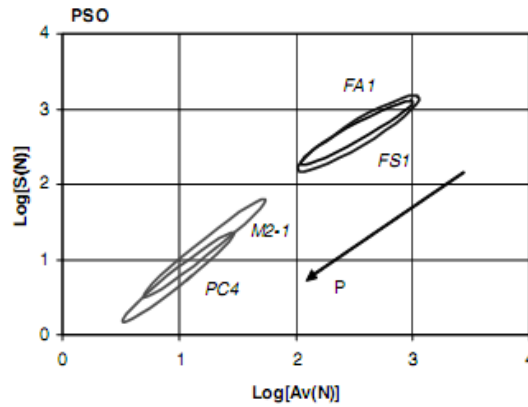


Fig. 3. $\text{Log}[S(N)]$ versus $\text{Log}[Av(N)]$ for $P = \{100, 500, 1000, 3000\}$ for the GA, the MA and the PSO algorithms, with $\delta = 0$.

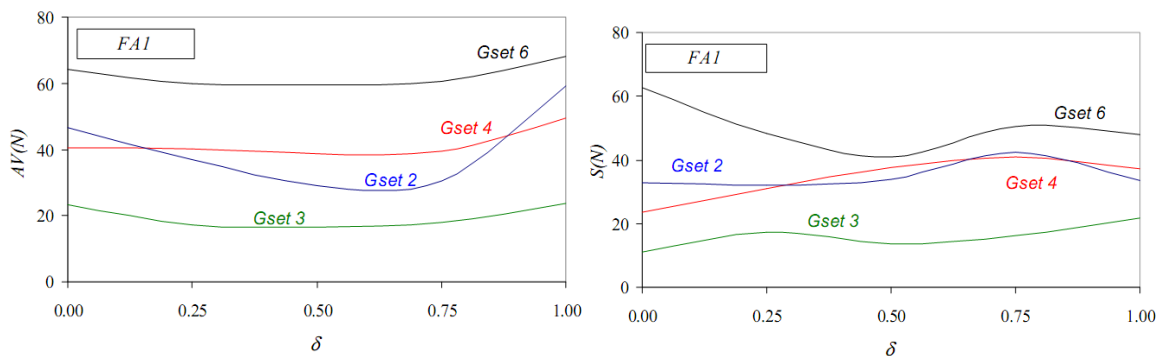


Fig. 4. Average number of generations $Av(N)$ and Standard deviation of the number of generations $S(N)$ to achieve a solution for the FA_1 circuit versus $\delta \in [0,1]$, with Gsets $\{6,4, 3, 2\}$.

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